

Detour to influences in Boolean FCNS.
and Talagrand's inequality (following Benjamini-Kalai-Schramm 2003)

Hypercube: $\{0,1\}^n$
Equipped with the uniform measure ρ .

Random variables: $f: \{0,1\}^n \rightarrow \mathbb{R}$.

Influence: For $1 \leq j \leq n$, let $\sigma_j f: \{0,1\}^n \rightarrow \mathbb{R}$

$$(\sigma_j f)(x) := f(x_1, \dots, x_{j-1}, 1-x_j, x_{j+1}, \dots, x_n)$$

Flip j 'th coordinate operation

$$(\rho_j f)(x) := \frac{f(x) - (\sigma_j f)(x)}{2}$$

The function measures the change made by the flip.
Norms of $\rho_j f$ can be thought of as the influence of the j 'th coordinate.

Norms: $\|f\|_p := (E\|f\|^p)^{1/p} = \left(\frac{1}{2^n} \sum_{x \in \{0,1\}^n} |f(x)|^p \right)^{1/p}$.

Theorem (Talagrand 1994):

$$\text{Var}(f) \leq C \sum_{j=1}^n \frac{\|\rho_j f\|_2^2}{1 + \log(\|\rho_j f\|_2 / \|\rho_j f\|_1)}$$

Without the denominator, this is in the same vein as the Efron-Stein inequality.

Get a significant improvement

if $\|\rho_j f\|_2 \gg \|\rho_j f\|_1$

This is the case, in our application, when $\rho_j f$ is only rarely non-zero.

To show a proof following

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Benjamini-Kalai-Schramm (2003) ¹⁹⁷⁰ ¹⁹⁷⁵

The proof relies on the Bonami-Beckner inequality

Bonami-Beckner: Noise operator N_ϵ ($0 \leq \epsilon \leq 1$) takes a function $f: \{0, 1\}^n \rightarrow \mathbb{R}$ and returns another fcn. which is somewhat averaged.

Let X_1, \dots, X_n be Bernoulli(ϵ) indep

$$X_i \sim \begin{cases} 1 & \epsilon \\ 0 & 1-\epsilon \end{cases}$$

$$N_\epsilon(f)(x) = \mathbb{E}(f(x + X \text{ mod } 2))$$

Each coord. is flipped with prob. ϵ and we average f on the resulting random position.

Theorem (Bonami-Beckner): For each $0 \leq \epsilon < 1$

$$\|N_\epsilon f\|_2 \leq \|f\|_1 + \delta(\epsilon)$$

Hypercontractive inequality

For $\delta(\epsilon) = (1-2\epsilon)^2$.

Representation in a Fourier-Walsh basis:

For a subset $S \subseteq \{1, \dots, n\} =: [n]$,

let $\chi_S: \{0, 1\}^n \rightarrow \{-1, 1\}$ be

$$\chi_S(x) = (-1)^{\sum_{i \in S} x_i}$$

Fourier-Walsh basis

$$(\chi_\emptyset \equiv 1)$$

This is an orthonormal basis for L^2 of fcn. on the hypercube with the uniform measure.

Parseval: Write $f(x) = \sum_S \hat{f}(S) \chi_S(x)$

OK

Parseval: Write $f(x) = \sum_{S \subseteq [n]} \hat{f}(S) \chi_S(x)$
 wrt. to the uniform measure \leftarrow coefficients in the expansion

$$\|f\|_2^2 = \sum_{S \subseteq [n]} |\hat{f}(S)|^2 \quad \leftarrow \text{wrt. the counting measure}$$

Noise operator: What is the expansion of $N_\epsilon f$?

Since $f(x) = \sum_{S \subseteq [n]} \hat{f}(S) \chi_S(x)$,

then $N_\epsilon(f)(x) = \sum_{S \subseteq [n]} \hat{f}(S) \mathbb{E}(\chi_S(x + X \text{ mod } 2))$

$$= \sum_{S \subseteq [n]} \hat{f}(S) \mathbb{E}((-1)^{\sum_{i \in S} x_i + X_i}) =$$

$$= \sum_{S \subseteq [n]} \hat{f}(S) \chi_S(x) \prod_{i \in S} \mathbb{E}((-1)^{X_i}) =$$

$= 1 - \epsilon - \epsilon = 1 - 2\epsilon$

$$= \sum_{S \subseteq [n]} (1 - 2\epsilon)^{|S|} \hat{f}(S) \chi_S(x).$$

Useful notation: For $-1 < p < 1$, write T_p

for the operator taking f to

$$T_p(f)(x) = \sum_{S \subseteq [n]} p^{|S|} \hat{f}(S) \chi_S(x)$$

So that $T_{1-2\epsilon} = N_\epsilon$.

In this notation, Bonami-Beckner says that

$$\|T_p\|_2 \leq \|f\|_{1+p^2}.$$

Proof of Talagrand's inequality:

Proof of Talagrand's inequality:

Write $f(x) = \sum_{S \subseteq [n]} \hat{f}(S) \chi_S(x)$.

Recall $(\rho_j f)(x) = \frac{f(x) - (\sigma_j f)(x)}{2}$ flip j'th input bit operation.

What is the Fourier-Walsh expansion of $\rho_j f$?

Simple to see that $(\sigma_j \chi_S)(x) = \begin{cases} \chi_S(x) & j \notin S \\ -\chi_S(x) & j \in S \end{cases}$

$\Rightarrow \rho_j(f)(x) = \sum_{\substack{S \subseteq [n] \\ j \in S}} \hat{f}(S) \chi_S(x)$

How to express the variance?

$\text{Var}(f) = \mathbb{E}(f - \mathbb{E}(f))^2 = \sum_{\phi \neq S \subseteq [n]} \hat{f}(S)^2$
Parseval

How to use the Bonami-Beckner ineq.?

$\text{Var}(f) = \sum_{\phi \neq S \subseteq [n]} \hat{f}(S)^2 = \sum_{j=1}^n \sum_{\phi \neq S \subseteq [n]} \frac{(\rho_j f)(S)^2}{|S|} \leq$

$\leq 3 \sum_{j=1}^n \int_0^1 \|T_p(\rho_j f)\|_2^2 dp$

Fourier-Walsh expansion = $\sum_{\substack{S \subseteq [n] \\ j \in S}} p^{|S|} \hat{f}(S) \chi_S(x)$

$\Rightarrow \|T_p(\rho_j f)\|_2^2 = \sum_{\substack{S \subseteq [n] \\ j \in S}} p^{2|S|} \hat{f}(S)^2$

Main inequality,

Bonami-Beckner $\Rightarrow \text{Var}(f) \leq 3 \sum_{j=1}^n \|T_p(\rho_j f)\|_2^2$.16

Main inequality,

$$\text{Bonami-Beckner} \Rightarrow \text{Var}(F) \leq 3 \sum_{j=1}^n \int_0^1 \| \rho_j F \|_{1+p^2}^2 dp.$$

Algebraic manipulations:

$$\begin{aligned} \mathbb{E}(|g|^{1+p^2}) &= \mathbb{E}[|g|^{2p^2} \cdot |g|^{1-p^2}] \leq (\mathbb{E}|g|^2)^{p^2} \cdot (\mathbb{E}|g|)^{1-p^2} \\ &= \|g\|_2^{2p^2} \cdot \|g\|_1^{1-p^2} = \|g\|_2^{1+p^2} \left(\frac{\|g\|_1}{\|g\|_2} \right)^{1-p^2}. \end{aligned}$$

$$\Rightarrow \|g\|_{1+p^2}^2 \leq \|g\|_2^2 \cdot \left(\frac{\|g\|_1}{\|g\|_2} \right)^{\frac{2(1-p^2)}{1+p^2}}.$$

Also note that for $0 \leq x \leq 1$:

$$\int_0^1 x^{\frac{2(1-p^2)}{1+p^2}} dp \leq \dots \leq \frac{1-x^{6/5}}{\log(\frac{1}{x})}.$$

Putting everything together:

$$\text{Var}(F) \leq 3 \sum_{j=1}^n \| \rho_j F \|_2^2 \cdot \frac{1 - (\| \rho_j F \|_1 / \| \rho_j F \|_2)^{6/5}}{\log(\| \rho_j F \|_2 / \| \rho_j F \|_1)}$$

one checks

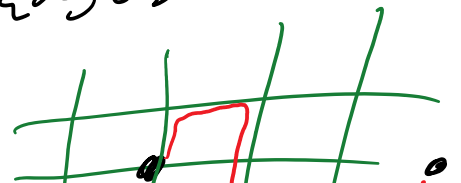
$$\leq C \sum_{j=1}^n \frac{\| \rho_j F \|_2^2}{1 + \log(\| \rho_j F \|_2 / \| \rho_j F \|_1)}$$

Finishing the proof of Talagrand's ineq.

Application to First Passage Percolation

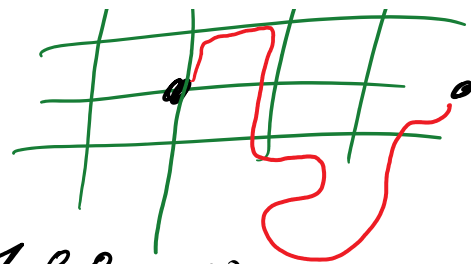
Consider FPP on \mathbb{Z}^d in which the edge weights are uniform on $\{a, b\}$

for some $0 < a < b < \infty$.



For some $0 < a < b < \dots$

Assume $d \geq 2$.



Consider $T(0, Le_1)$ ($e_1 = (1, 0, 0, \dots, 0)$, first coord. dir.)
which is the shortest passage time from 0 to Le_1 .

Thm. (Benjamini-Kalai-Schramm 2003):

$$\text{Var}(T(0, Le_1)) \leq C \frac{L}{\log L}.$$

(Best known upper bound)
(Believed in $d=2$ that $\text{var} \approx L^{2/3}$
and even smaller in $d \geq 3$).

Note that the optimal path can use at most $\frac{b}{a} \cdot L$ edges since the straight line path is a candidate.

Thus $T(0, Le_1)$ is a fcn. of finitely many Boolean random variables. So Talagrand's Ineq. applies.

For brevity, write $F = T(0, Le_1)$,
a fcn. of (η_e) , where $\eta_e \in \{a, b\}$ is the weight of the edge e .

Note $\rho_e F = \frac{F - \sigma_e F}{2}$, where $\sigma_e F$ is the passage time after flipping the weight of the edge e .

Note: if $\eta_e = a$ then $\sigma_e F \geq F$ always

Note: if $\eta_e = a$ then $\sigma_e F \geq F$ always
 and $\sigma_e F > F$ iff all shortest paths
 from o to Le_1 pass through e .

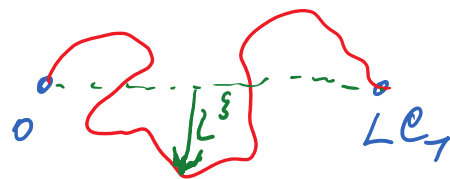
$$\text{Thus } \mathbb{E}(\rho_e F)^2 = \frac{1}{2} \mathbb{E}((\rho_e F)^2 | \eta_e = a) =$$

$$= \mathbb{P}(\text{all optimal paths} | \eta_e = a) \cdot c(a, b).$$

To improve upon Efron-Stein, need
 that $\frac{\|\rho_e F\|_2}{\|\rho_e F\|_1}$ is large.

$$\text{But } \frac{\|\rho_e F\|_2}{\|\rho_e F\|_1} = \frac{c(a, b)}{\sqrt{\mathbb{P}(\text{all optimal paths} | \eta_e = a)}}$$

This probability is expected to be
 small for edges far from o and Le_1 ,
 probably as small as $\frac{1}{L^c}$,
 and this would imply the BKS
 upper bound on the variance,
 by Talagrand's ineq.



Open question: Prove such an upper bound
 on the probability (BKS midpoint problem).

Known: The above prob. tends to 0 with L
 for e far from o and Le_1 .

(Ahlberg-Hoffman 2016)

Idea to bypass the open question:

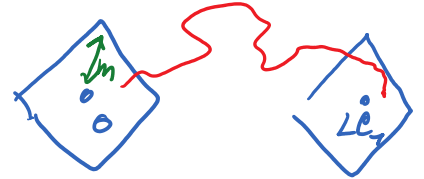
Idea to bypass the open question:

Instead of $T(0, Le_1)$ consider an averaged

version:
$$S(0, Le_1) := \frac{1}{|B(0, m)|} \sum_{z \in B(0, m)} T(z, z + Le_1)$$

where $B(0, m) := \{z \in \mathbb{Z}^d : \|z\|_1 \leq m\}$,
and choosing $m := L^{1/4}$.

Certainly, $\mathbb{E} S(0, Le_1) = \mathbb{E} T(0, Le_1)$
by translation invariance.



Additionally, $|S(0, Le_1) - T(0, Le_1)| \leq \underbrace{C(a, b)}_{2b} \cdot m$

$$\Rightarrow \text{Var}(T) = \|T - \mathbb{E}T\|_2^2 \leq \underbrace{2\|S - \mathbb{E}S\|_2^2}_{=\text{Var}(S)} + 2(C(a, b)m)^2.$$

\Rightarrow It suffices to use Talagrand's ineq. for S . The advantage is that for S , due to the averaging, one can show that $\|P_{e_1} S\|_2 / \|P_{e_1} S\|_1$ is a power of L^{-1} without much difficulty (see 50 years of FPP survey).

Final remarks: The upper bound $\text{Var}(T) \leq \frac{dL}{\log L}$ is now known for most weight dist.

Talagrand's ineq. is sometimes replaced by Falik-Smorodnitsky's ineq.

Damron-Hanson-Sasoe 2015

no four to total variation distances

Detour to total variation distances and a lower bound on the variance

Best-known lower bound on the variance,

Thm. (Newman-Piza 1995): In $d=2$,

$$\text{Var}(T(L, L_{\uparrow})) \geq c \log L.$$

open whether $\text{Var} \xrightarrow{L \rightarrow \infty} \infty$ in $d \geq 3$

(open even in physics literature).

← in $d=2$,
expect
 $\text{Var} \approx L^{2/3}$.

Total variation distance:

μ, ν are prob. measures on some meas. space.

$$d_{TV}(\mu, \nu) := \sup_{A \text{ event}} |\mu(A) - \nu(A)|$$

For every A , $\nu(A) - d_{TV}(\mu, \nu) \leq \mu(A) \leq \nu(A) + d_{TV}(\mu, \nu)$

When μ, ν have densities wrt. to a third measure λ (always ok with, say, $\lambda = \mu + \nu$) (usually, $\lambda = \text{Lebesgue or counting}$)

write $\mu = f d\lambda$, $\nu = g d\lambda$.

Claim: $d_{TV}(\mu, \nu) = \frac{1}{2} \int |f - g| d\lambda =$

$$= \int (f - g) \mathbb{1}_{f > g} d\lambda = \int (g - f) \mathbb{1}_{g > f} d\lambda.$$

Proof: First, $\frac{1}{2} \int |f - g| d\lambda =$

$$= \frac{1}{2} \left[\int (f - g) \mathbb{1}_{f > g} d\lambda + \int (g - f) \mathbb{1}_{g > f} d\lambda \right] =$$

$$= \frac{1}{2} \left[\int (F-g) \mathbb{I}_{F>g} d\lambda + \underbrace{\int (g-F) d\lambda}_{=0} - \int (g-F) \mathbb{I}_{F<g} d\lambda \right]$$

$$= \int (F-g) \mathbb{I}_{F>g} d\lambda.$$

Now, for any event A ,

$$\begin{aligned} \mu(A) &= \int_A F d\lambda \leq \int_A g d\lambda + \int_A (F-g) \mathbb{I}_{F>g} d\lambda \leq \\ &\leq \nu(A) + \int (F-g) \mathbb{I}_{F>g} d\lambda \end{aligned}$$

$$\Rightarrow d_{TV}(\mu, \nu) \leq \int (F-g) \mathbb{I}_{F>g} d\lambda.$$

For other direction,

if $A = \{F > g\}$ then

$$\mu(A) = \int_A F d\lambda = \int_A g d\lambda + \int (F-g) \mathbb{I}_{F>g} d\lambda.$$

$\underbrace{\int_A g d\lambda}_{= \nu(A)}$

It is not always easy to give an upper bound to $d_{TV}(\mu, \nu)$. We discuss now one of the simplest cases.

Example: $\underline{X} = (X_1, \dots, X_n)$ IID, $X_i \sim N(0,1)$

and $\underline{Y} = (Y_1, \dots, Y_n)$ indep.

where $Y_i \sim N(\epsilon_i, 1)$.

How big is $d_{TV}(\mathcal{L}(\underline{X}), \mathcal{L}(\underline{Y}))$? distributed as $N(0,1)$

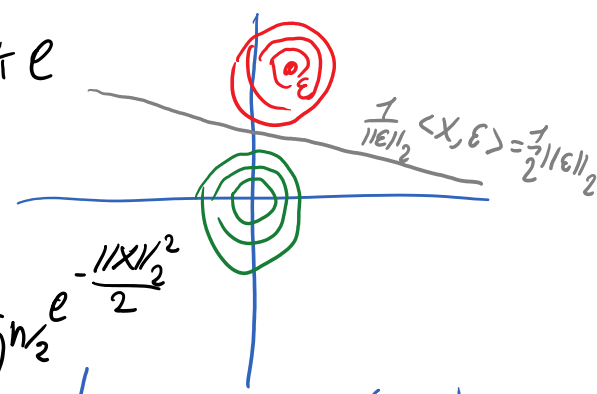
Claim: $d_{TV}(\mathcal{L}(\underline{X}), \mathcal{L}(\underline{Y})) = \mathbb{P}\left(\left| \frac{1}{\|\epsilon\|_2} \langle \underline{X}, \epsilon \rangle \right| < \frac{1}{2} \|\epsilon\|_2\right)$

Claim: $d_{TV}(\mathcal{L}(X), \mathcal{L}(Y)) = \mathbb{P}\left(\left|\frac{1}{\|E\|_2} \langle X, E \rangle\right| < \frac{1}{2} \|E\|_2\right)$

$\approx \begin{cases} c \|E\|_2 & \|E\|_2 \ll 1 \\ 1 - \frac{c}{\|E\|_2} e^{-\frac{1}{8} \|E\|_2^2} & \|E\|_2 \gg 1 \end{cases}$

Proof by picture: densities for X and Y

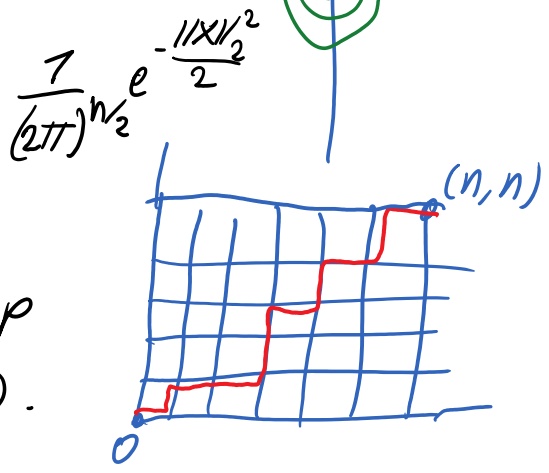
wrt. Leb. measure bok \mathbb{R}^k



Application to directed last passage percolation with Gaussian weights

$T(0, n) =$ maximal sum of weights over all right-up paths from $(0, 0)$ to (n, n) .

Here we take the weights to be IID $N(0, 1)$.



Claim: $\text{Var}(T(0, n)) \geq c \text{Log} n$.

Proof: Let (η_e) be the edge weights. Define new weights $\tilde{\eta}_e = \eta_e + \delta_e$, where (δ_e) are deterministic and depending only on $|e| :=$ distance of e from $(0, 0)$. Define $\tilde{T}(0, n)$ to be like $T(0, n)$ but with the weights $(\tilde{\eta}_e)$.

Clearly, $\tilde{T}(0, n) \geq T(0, n) + \sum_{|e|=1}^{2n} \delta_{|e|}$

Since the old geodesic is a candidate for the optimal path (and actually it equals the maximal path)

Since the ...
the optimal path (and actually it equals the optimal path).

Now consider the event

$$A = \{ |T(0, n) - \mathbb{E}(T(0, n))| \leq t \}$$

Then $|P(\eta \in A) - P(\tilde{\eta} \in A)| \leq d_{TV}(\mathcal{L}(\eta), \mathcal{L}(\tilde{\eta}))$.

If we take t to be, say, two standard deviations of $T(0, n)$ then by Chebyshev's inequality, $P(\eta \in A) \geq \frac{3}{4}$.

Thus, if $d_{TV}(\mathcal{L}(\eta), \mathcal{L}(\tilde{\eta})) \leq \frac{1}{2}$, say, then A is likely also for $\tilde{\eta}$,

whence $\sum_{|e|=1}^{2n} \delta_{|e|} \leq 2t =$ four standard deviations of $T(0, n)$.

Conclude: $\text{std}(T(0, n)) \geq c \sum_{|e|=1}^{2n} \delta_{|e|}$

Whenever $d_{TV}(\mathcal{L}(\eta), \mathcal{L}(\tilde{\eta})) \leq \frac{1}{2}$.

Choice of δ_e : choose $\delta_{|e|} = \frac{c}{|e| \sqrt{\log n}}$.

Then $d_{TV}(\mathcal{L}(\eta), \mathcal{L}(\tilde{\eta})) \approx \|\delta_e\|_2 = \sum_{|e|=1}^{2n} \delta_{|e|}^2 \cdot |e| \leq \frac{1}{2}$.

$\Rightarrow \text{std}(T(0, n)) \geq c \sum_{|e|=1}^{2n} \delta_{|e|} = c \sqrt{\log n}$.

Mermin-Wagner style arguments